

# XcalableMP and XcalableACC for Productivity and Performance in HPC Challenge Award Competition

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# Outline

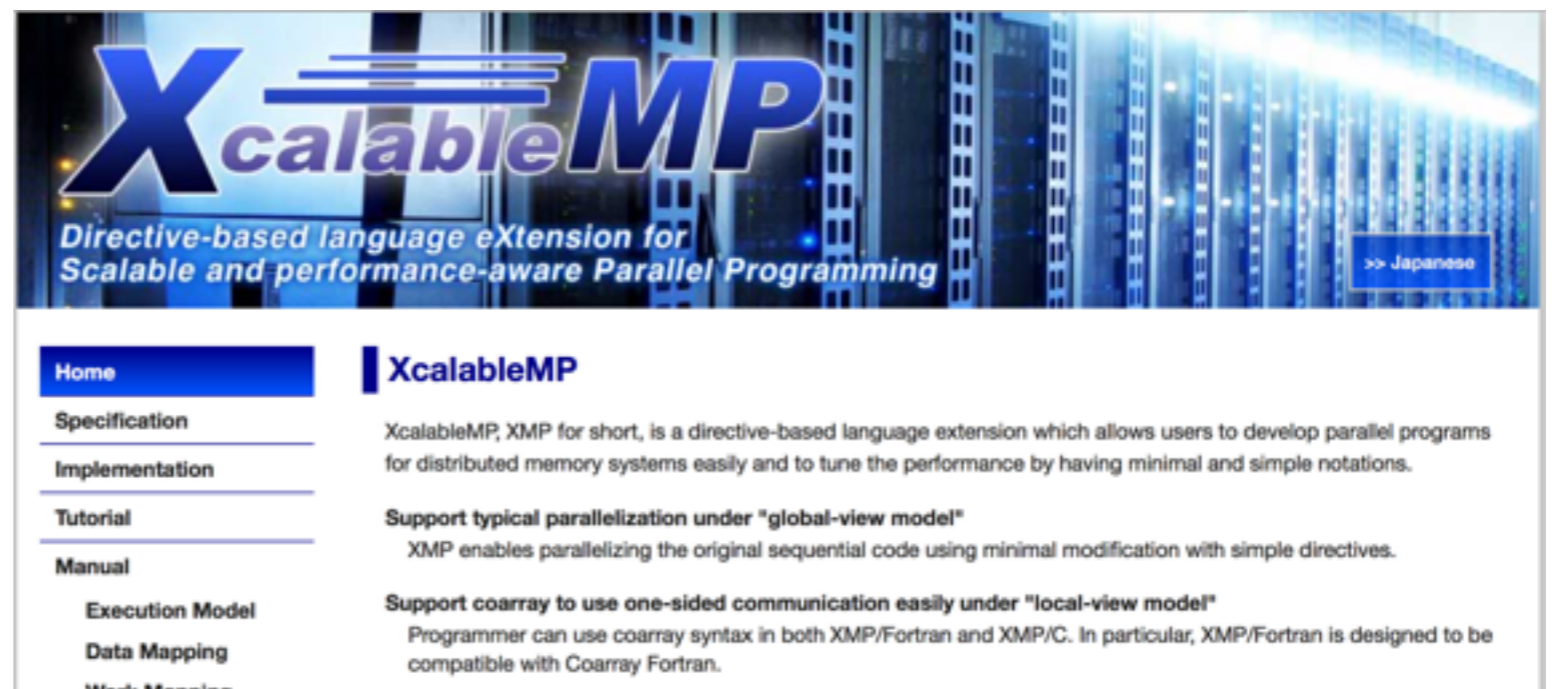
1. XcalableMP (XMP) for cluster systems (14min.)

2. XcalableACC (XACC) for accelerator cluster systems (6min.)

Extension of XMP using OpenACC

Sorry !!, work-in-progress

The submission report is available at <http://xcalablemp.org>



**XcalableMP**  
Directive-based language extension for Scalable and performance-aware Parallel Programming

>> Japanese

- Home
- Specification
- Implementation
- Tutorial
- Manual
  - Execution Model
  - Data Mapping
  - Work Mapping

### XcalableMP

XcalableMP, XMP for short, is a directive-based language extension which allows users to develop parallel programs for distributed memory systems easily and to tune the performance by having minimal and simple notations.

**Support typical parallelization under "global-view model"**  
XMP enables parallelizing the original sequential code using minimal modification with simple directives.

**Support coarray to use one-sided communication easily under "local-view model"**  
Programmer can use coarray syntax in both XMP/Fortran and XMP/C. In particular, XMP/Fortran is designed to be compatible with Coarray Fortran.

# What is XcalableMP (XMP) ?

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## Directive-based language extensions of Fortran and C

- By XMP specification working group of PC cluster consortium (SC Booth#2924)
- Version 1.2.1 specification available (<http://xcalablemp.org>)

## Support two memory models

- **Global-view** (HPF-like data/work mapping directives)
- **Local-view** (coarray)

## Implementation of Compiler

- Omni XMP Compiler version 0.9 (<http://omni-compiler.org>)
- Platforms: Fujitsu the K computer and FX10, Cray XT/XE, IBM BlueGene, NEC SX, Hitachi SR, Linux clusters, etc.

# Code example (Global-view)

---

```
int a[MAX];
```

```
#pragma xmp nodes p(4)
```

```
#pragma xmp template t(0:MAX-1)
```

```
#pragma xmp distribute t(block) on p
```

Data distribution

```
#pragma xmp align a[i] with t(i)
```

```
main(){
```

```
  int i, j, res = 0;
```

add to the serial code : incremental parallelization

```
#pragma xmp loop on t(i) reduction(+:res)
```

```
  for(i = 0; i <MAX; i++){
```

```
    a[i] = func(i);
```

```
    res += array[i];
```

```
  }
```

Work mapping and data synchronization

# Code example (Local-view)

---

```
double a[100]:[*], b[100]:[*];  
int me = xmp_node_num();
```

Define Coarrays

```
if(me == 2)  
    a[:,1] = b[:,];
```

Put Operation

```
if(me == 1)  
    a[0:50] = b[0:50]:[2];
```

Get Operation

## Coarray syntax in XMP/C

```
array_name[start:length]:[node_number];
```

**XMP/Fortran is upward compatible with Fortran 2008**

# Results and Machine

## Summary

Four HPCC Benchmarks

Benchmark		# Nodes	Performance (/peak)	SLOC
HPL	Ver. 1	16,384	971 TFlops (46.3%)	313
	Ver. 2	4,096	423 TFlops ( <b>80.7%</b> )	426
FFT		<b>82,944</b>	212 TFlops (2.0%)	205
STREAM		<b>82,994</b>	3,583 TB/s (67.5%)	69
RandomAccess		16,384	254 GUPs	253

## The K computer: 82,944 nodes



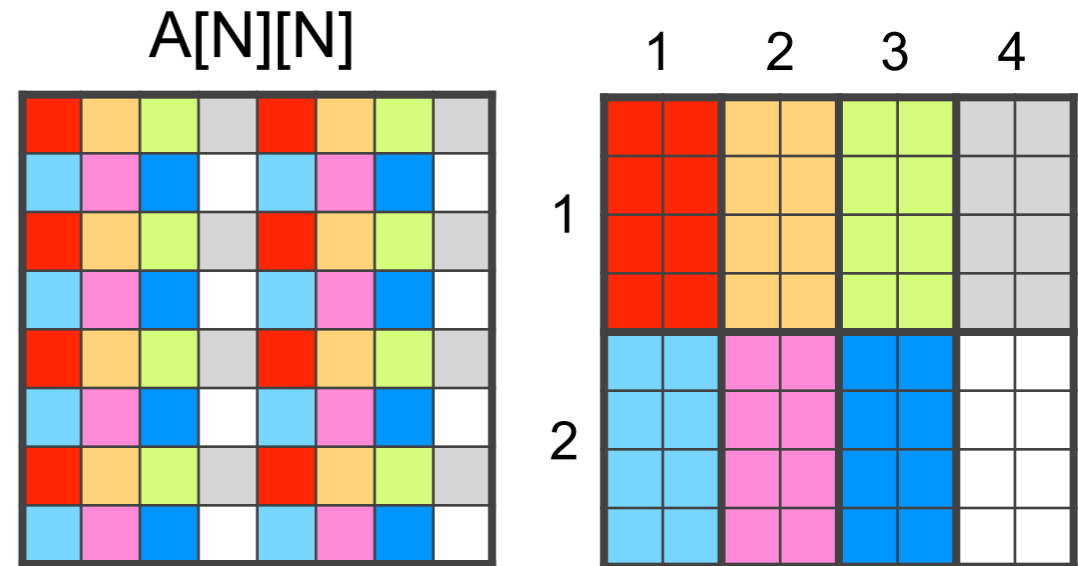
<http://www.aics.riken.jp/jp/outreach/photogallery.html>

- SPARC64 VIIIfx Chip, 128 GFlops
- DDR3 SDRAM 16GB, 64GB/s
- Tofu Interconnect
  - 6D mesh/torus network
  - 5GB/s x 4links x 2

# HPL version 1

- Source lines of Code (SLOC) is **313**, written in XMP/C
- Block-Cyclic Distribution

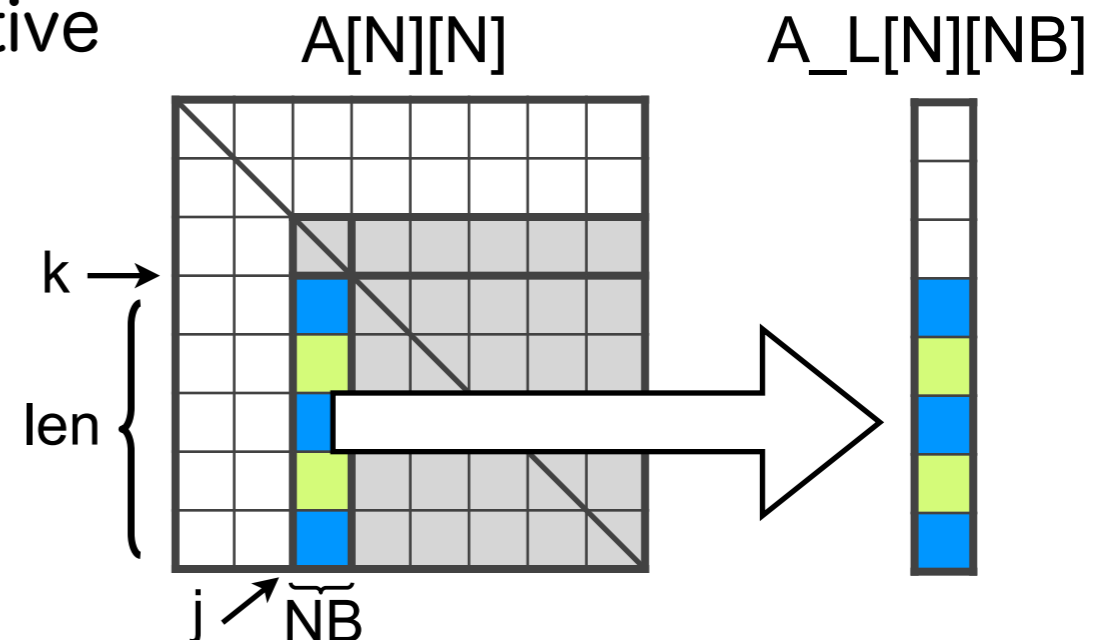
```
double A[N][N];
#pragma xmp nodes p(P,Q)
#pragma xmp template t(0:N-1, 0:N-1)
#pragma xmp distribute t(cyclic(NB), \
                        cyclic(NB)) onto p
#pragma xmp align A[i][j] with t(j,i)
```



Programmer can use BLAS for distributed array.

- Panel Broadcast by using **gmove** directive

```
double A_L[N][NB];
#pragma xmp align A_L[i][*] with t(*,i)
:
#pragma xmp gmove
A_L[k:len][0:NB] = A[k:len][j:NB];
```

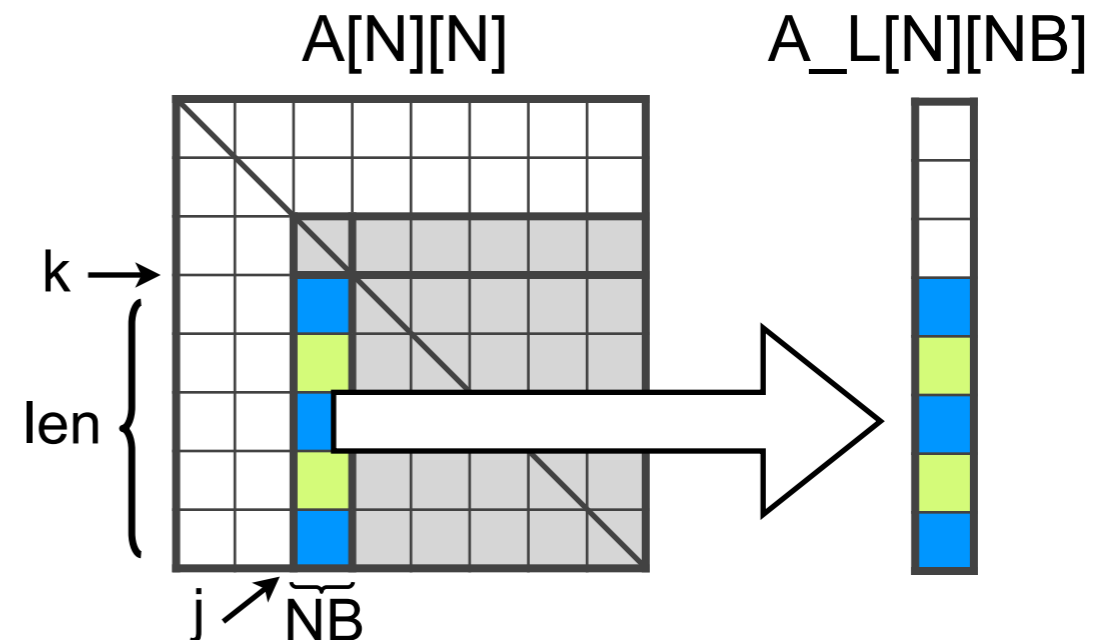


# HPL version 2

- SLOC is **426**, written in XMP/C
- "Lookahead algorithm" by using **gmove** directive with **async** clause

Overlap communication and calculation

```
double A_L[N][NB];  
#pragma xmp align A_L[i][*] with t(*,i)  
:  
#pragma xmp gmove async(1)  
A_L[k:len][0:NB] = A[k:len][j:NB];  
:  
for(m=j+NB;m<N;m+=NB){  
  for(n=j+NB;n<N;n+=NB){  
    cblas_dgemm(&A[m][n], ..);  
    if(xmp_test_async(1)){  
      // receive A[k:len][j:NB];  
      :  
    }  
  }  
}
```

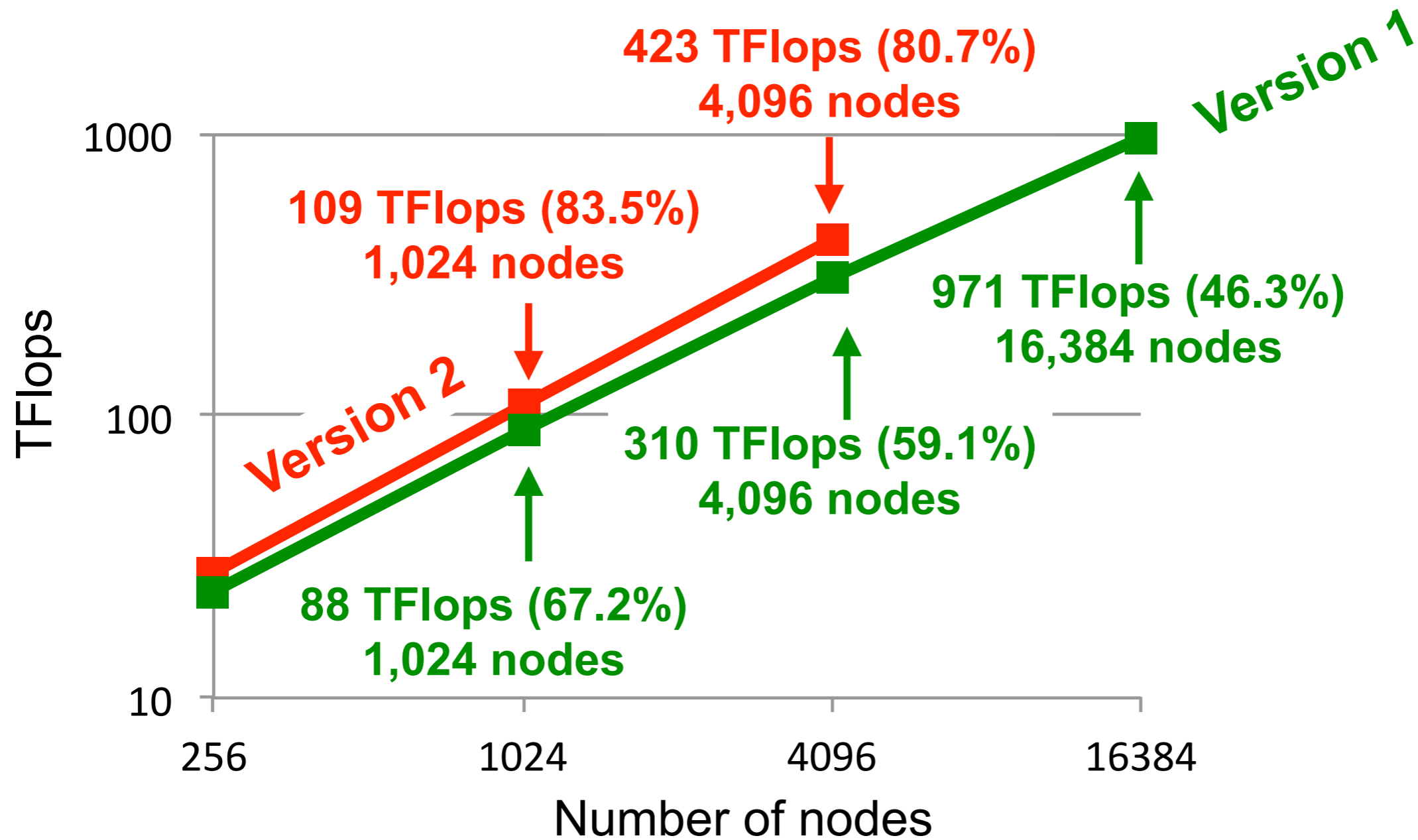


asynchronous broadcast  
communication

**Confirm whether data with async clause comes or not.**



# Performance of HPL



XMP-HPL Version 2 has a good scalability.

Sorry, the measurement in 16,384 nodes is late for this BoF.

# RandomAccess

- SLOC is **253**, written in XMP/C
- Local-view programming with XMP/C coarray syntax
- The XMP RandomAccess is iterated over sets of CHUNK updates on each node

```
u64Int recv[LOGPROCS][RCHUNK+1]:[*];
```

```
...
```

```
for (j = 0; j < logNumProcs; j++) {  
    recv[j][0:num]:[i_partner] = send[i][0:num];
```

```
#pragma xmp sync_memory
```

```
#pragma xmp post(p(i_partner), 0)
```

```
:
```

```
#pragma xmp wait(p(j_partner))
```

```
}
```

← Define coarray

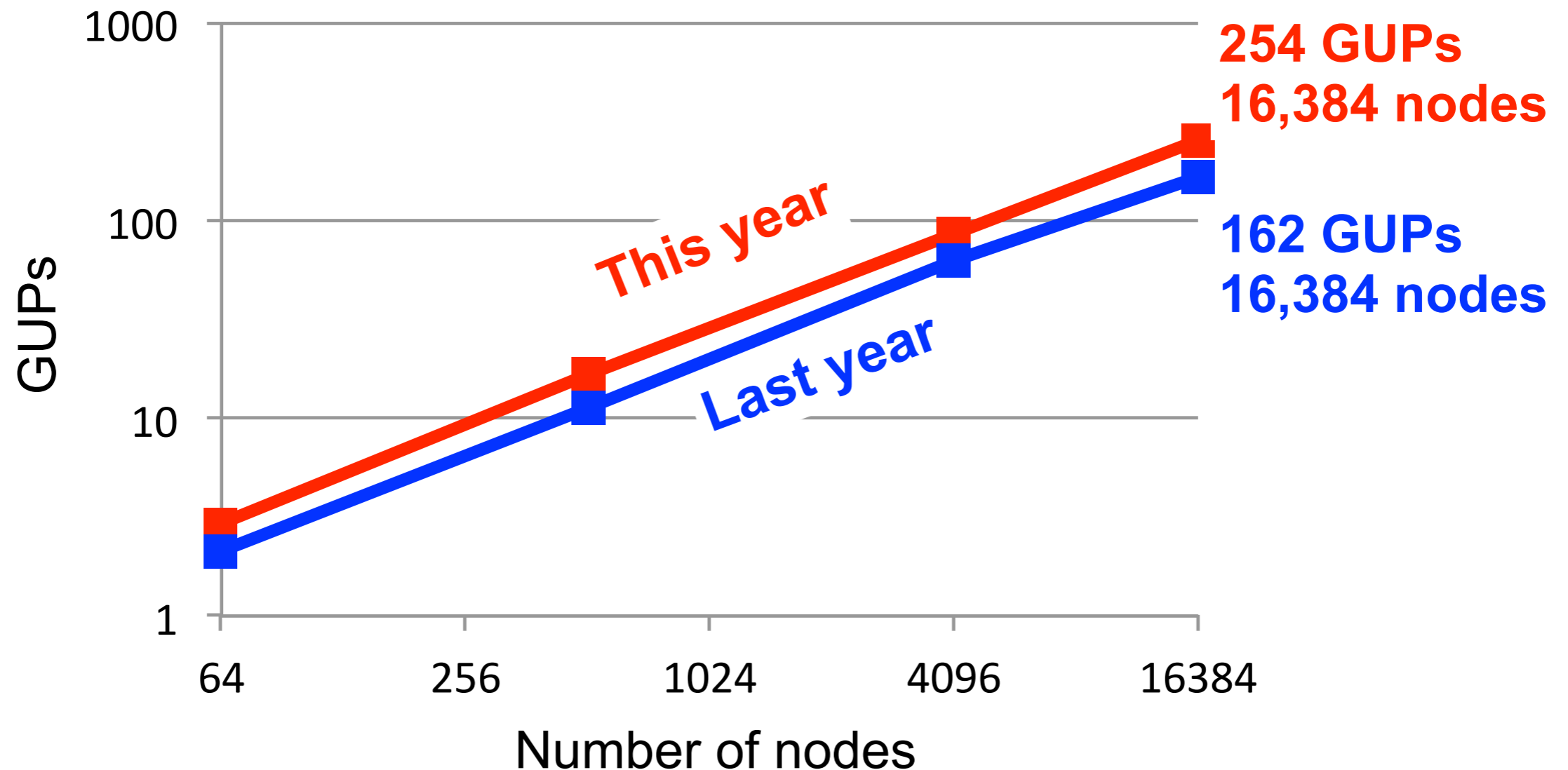
← Put operation

←

A point-to-point synchronization is specified with the XMP's post and wait directives to realize asynchronous behavior of this algorithm

# Performance of RandomAccess

Last year, to implement the post/wait directives, XMP uses MPI\_Send/Recv.  
This year, to implement them, XMP uses RDMA of the K computer.

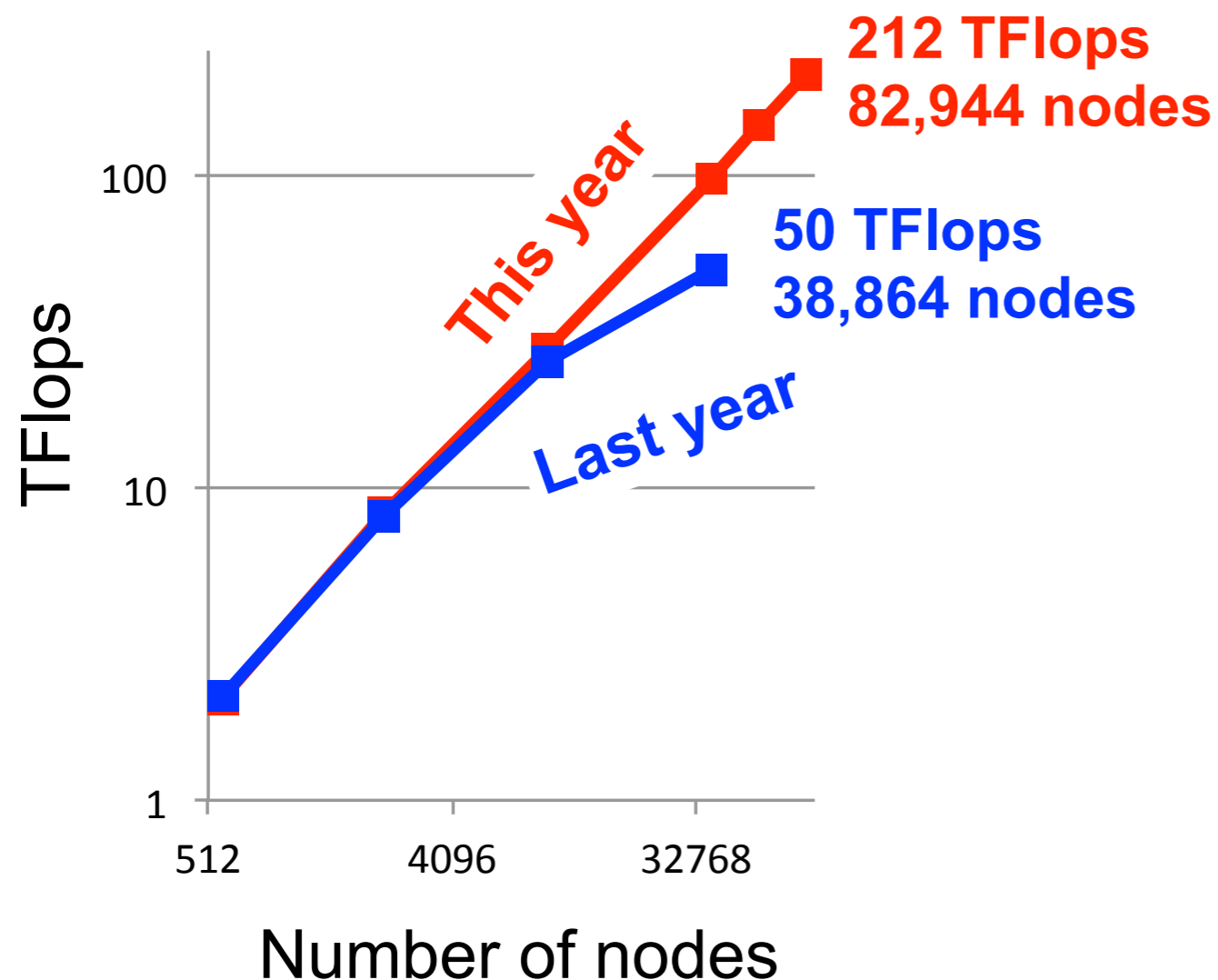


# FFT and STREAM

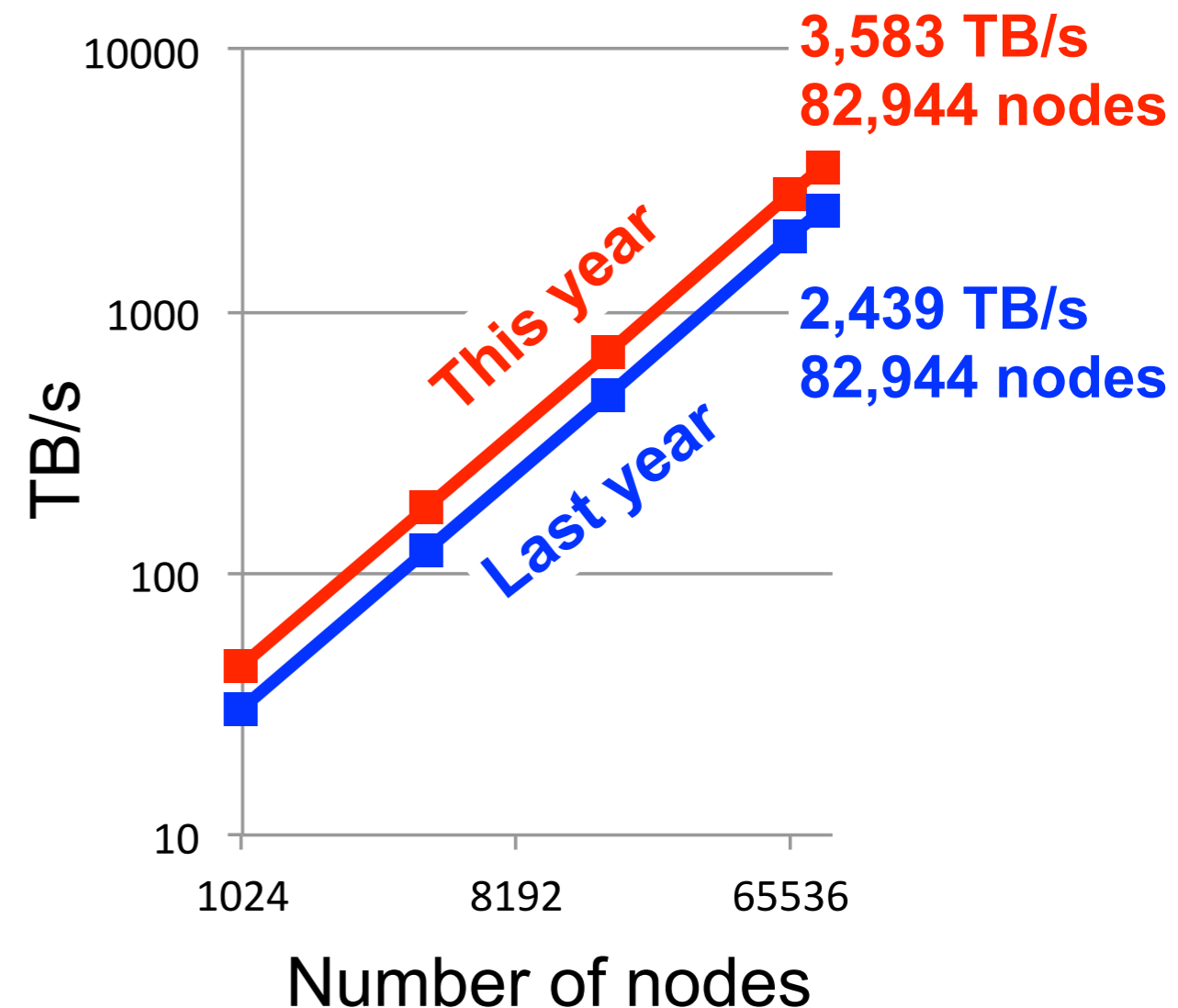
Code cleanup and performance improvement.

Please refer to the submission report at <http://xcalablemp.org>

## FFT (SLOC 205, XMP/F)

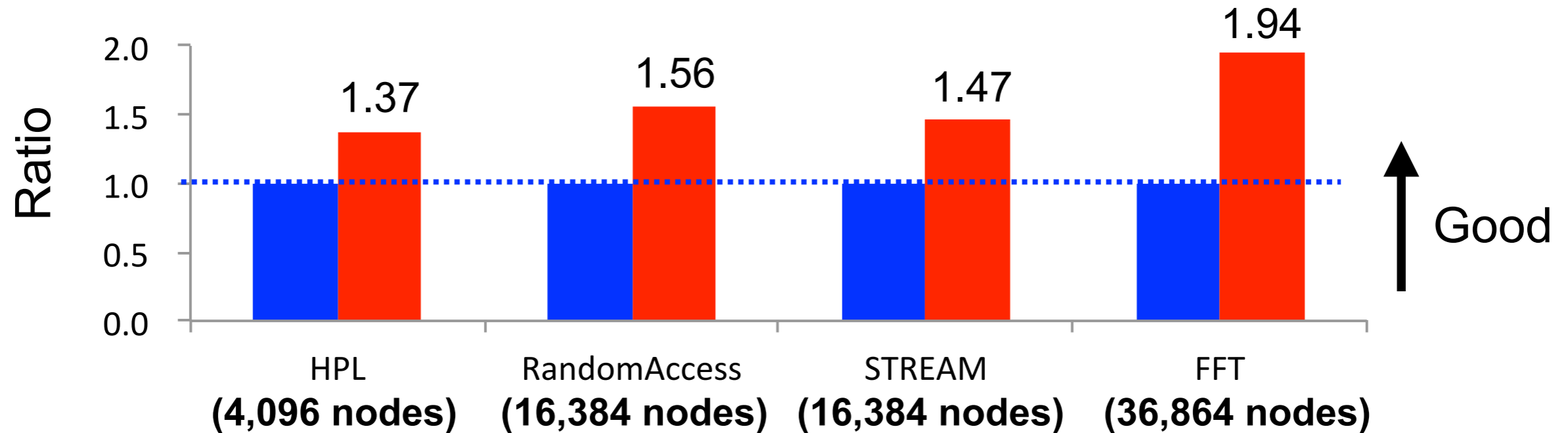


## STREAM (SLOC 69, XMP/C)

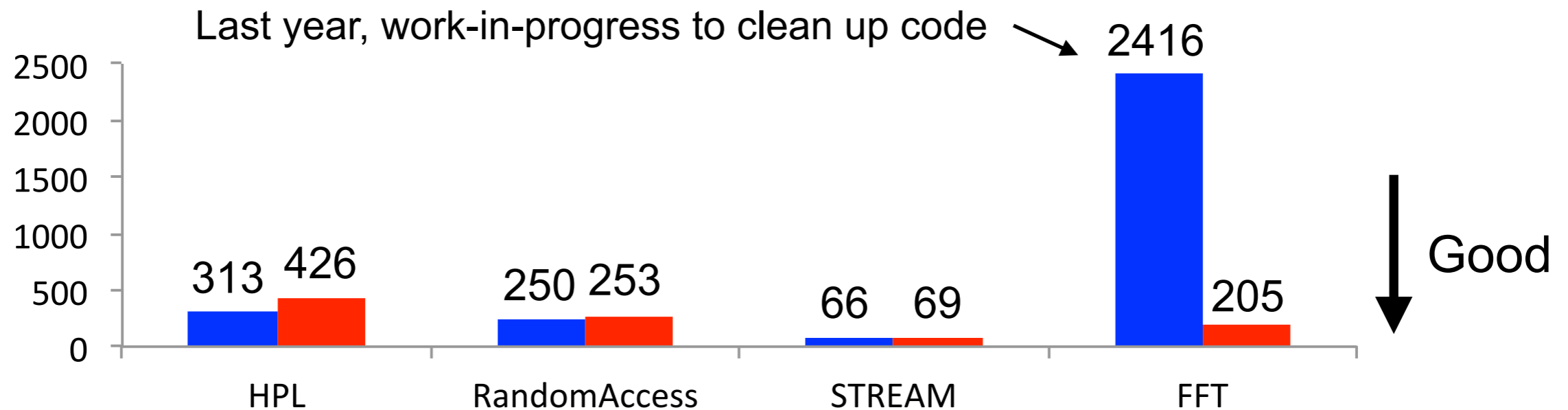


# Compare to two versions

- Improvement rate (on the same nodes) **37 - 94% improvement !!**



- SLOC



# Outline

1. XcalableMP (XMP) for cluster systems

(14min.)

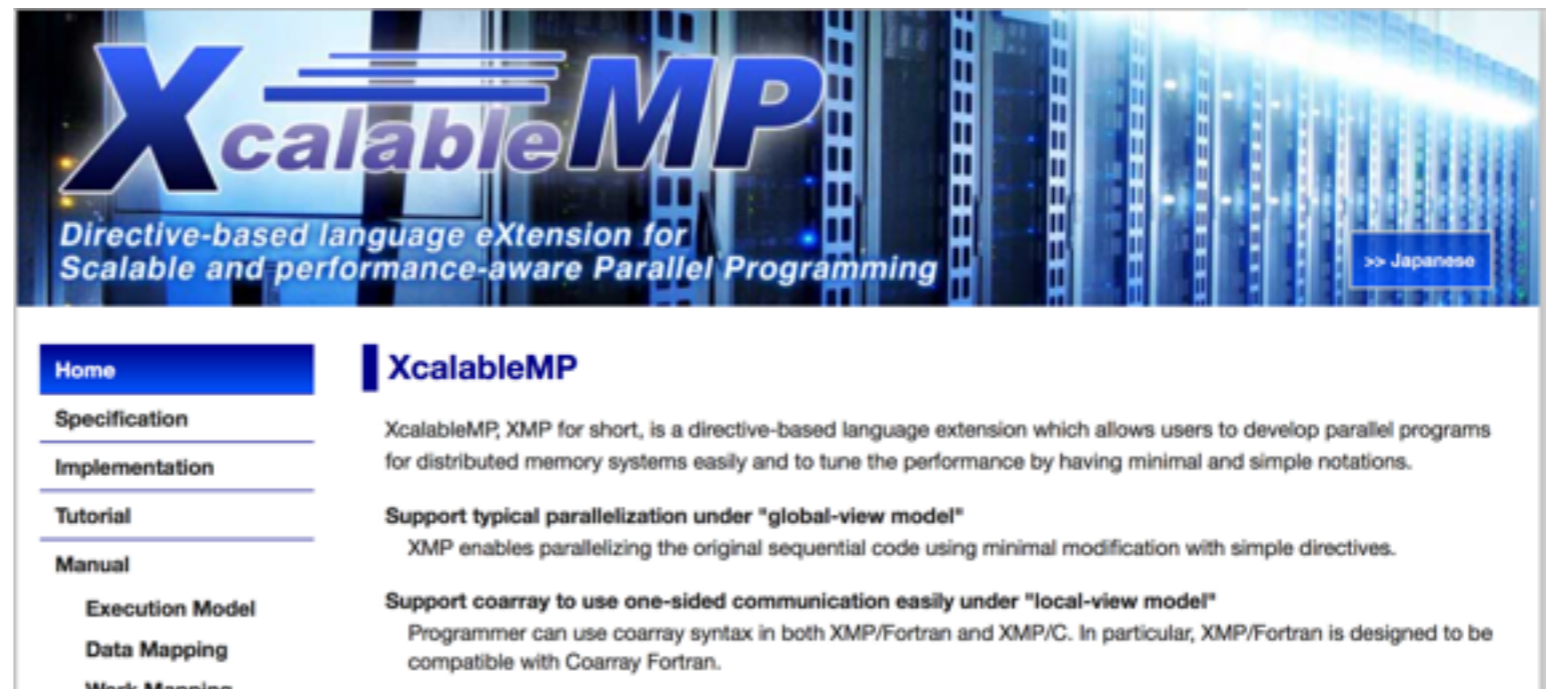
2. XcalableACC (XACC) for accelerator cluster systems

(6min.)

Extension of XMP using OpenACC

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# What is XcalableACC?

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**Extension of XMP using OpenACC for accelerator clusters**

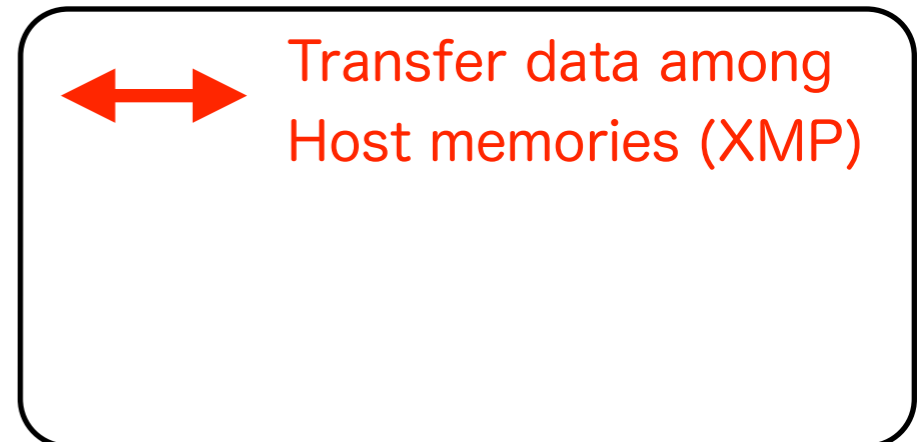
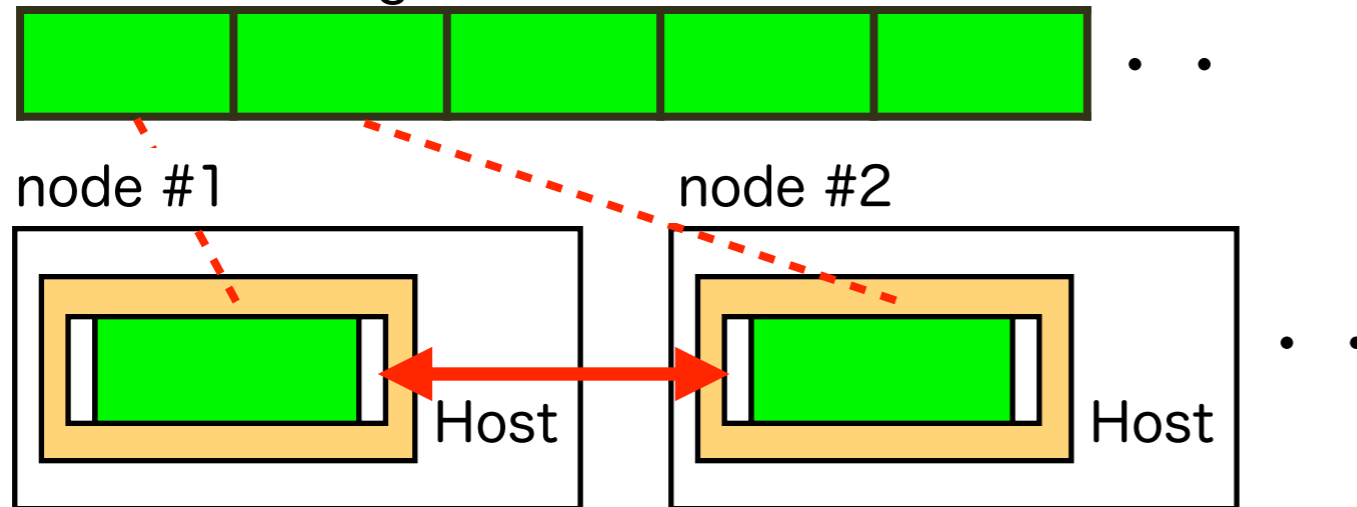
## **Feature:**

- Mix XMP and OpenACC directives seamlessly
- Support transferring data among accelerators directly

# Difference XMP and XACC memory models

- **XMP** memory model

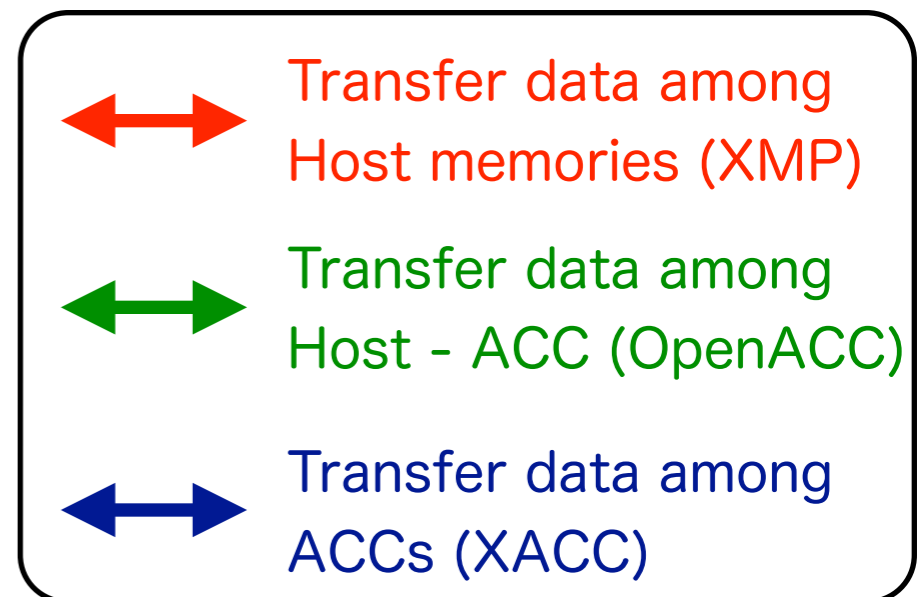
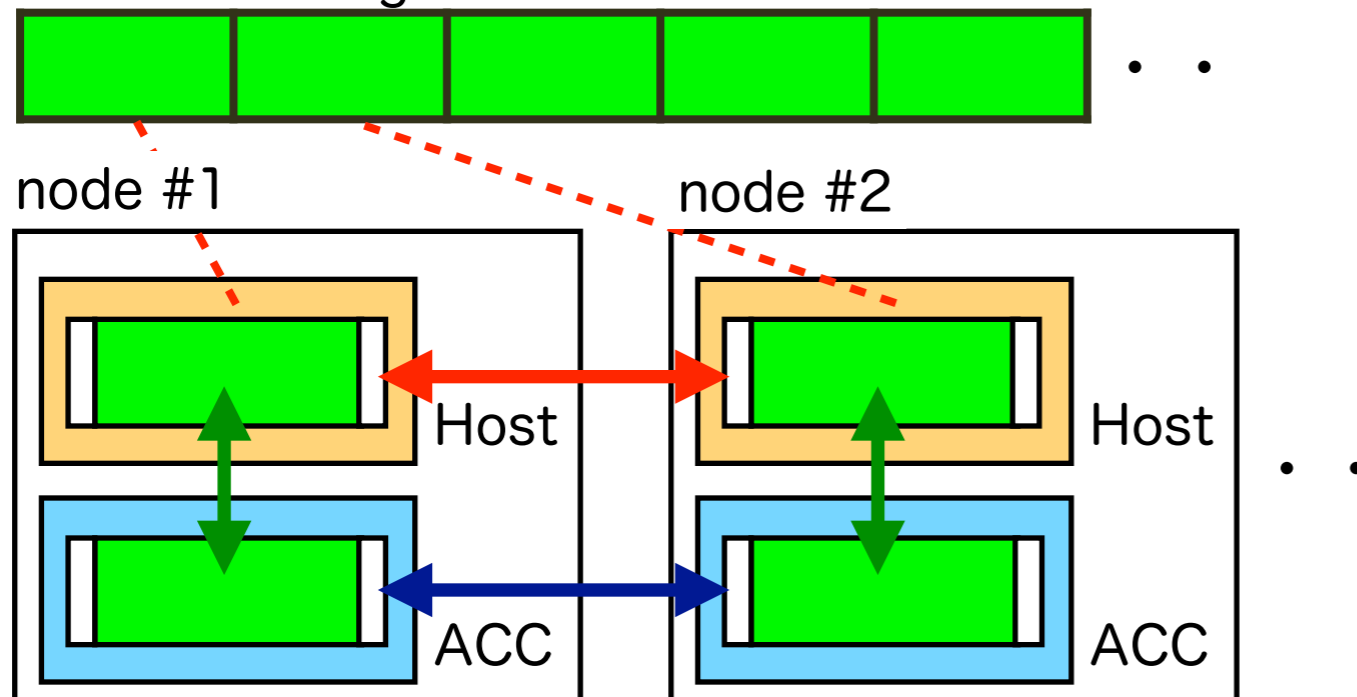
Global Indexing



- **XACC** memory model

Map “global Indexing” to accelerators

Global Indexing





# XACC code example

```
double u[XSIZE][YSIZE], uu[XSIZE][YSIZE];
#pragma xmp nodes p(x, y)
#pragma xmp template t(0:YSIZE-1, 0:XSIZE-1)
#pragma xmp distribute t(block, block) onto p
#pragma xmp align [j][i] with t(i,j) :: u, uu
#pragma xmp shadow uu[1:1][1:1]
...
#pragma acc data copy(u) copyin(uu)
{
  for(k=0; k<MAX_ITER; k++){
    #pragma xmp loop (y,x) on t(y,x)
    #pragma acc parallel loop collapse(2)
      for(x=1; x<XSIZE-1; x++)
        for(y=1; y<YSIZE-1; y++)
          uu[x][y] = u[x][y];

    #pragma xmp reflect (uu) acc

    #pragma xmp loop (y,x) on t(y,x)
    #pragma acc parallel loop collapse(2)
      for(x=1; x<XSIZE-1; x++)
        for(y=1; y<YSIZE-1; y++)
          u[x][y] = (uu[x-1][y]+uu[x+1][y]+
                    uu[x][y-1]+uu[x][y+1])/4.0;
  } // end k
} // end data
```

## Laplace's equation

Data Distribution

Transfer XMP distributed arrays to accelerator

OpenACC directive parallelizes the loop statement parallelized by XMP directive

Exchange halo region of uu[][]

When “acc” clause is specified in XMP communication directive, data on accelerator is transferred.

# Results and Machine

## Summary

Three HPCC Benchmarks and HIMENO Benchmark

Benchmark	#Nodes	#CPUs	#GPUs	Performance (/peak)	SLOC
HPL	32	64	128	7 TFlops (4.2%)	343
FFT	32	64	-	257 GFlops (0.1%)	205
STREAM	64	128	256	15 TB/s (20.4%)	84
HIMENO	64	128	256	14 TFlops (1.4%)	253

## HA-PACS/TCA: 64 nodes



- Ivy Bridge E5-2680v2, 224GFlops x 2 Sockets
- DDR3 SDRAM 128GB, 59.7GB/s x 2
- Infiniband 4xQDR x 2 rails : 8GB/s
- **NVIDIA K20X (4GPUs / Node)**
  - 1.31 TFlops/GPU(SP), 3.95 TFlops/GPU(DP)
  - 250GB/s/GPU

<http://www.ccs.tsukuba.ac.jp/CCS/eng/research-activities/projects/ha-pacs>

# STREAM

The XACC STREAM uses both CPUs and GPUs together, **XMP**, **OpenACC**, and **OpenMP** directives are used.

```
#pragma xmp nodes p(*)
#pragma acc data copy(a[0:GPU_SIZE], b[0:GPU_SIZE], c[0:GPU_SIZE])
{
  for(k=0; k<NTIMES; k++) {
#pragma xmp barrier
    times[k] = -xmp_wtime();

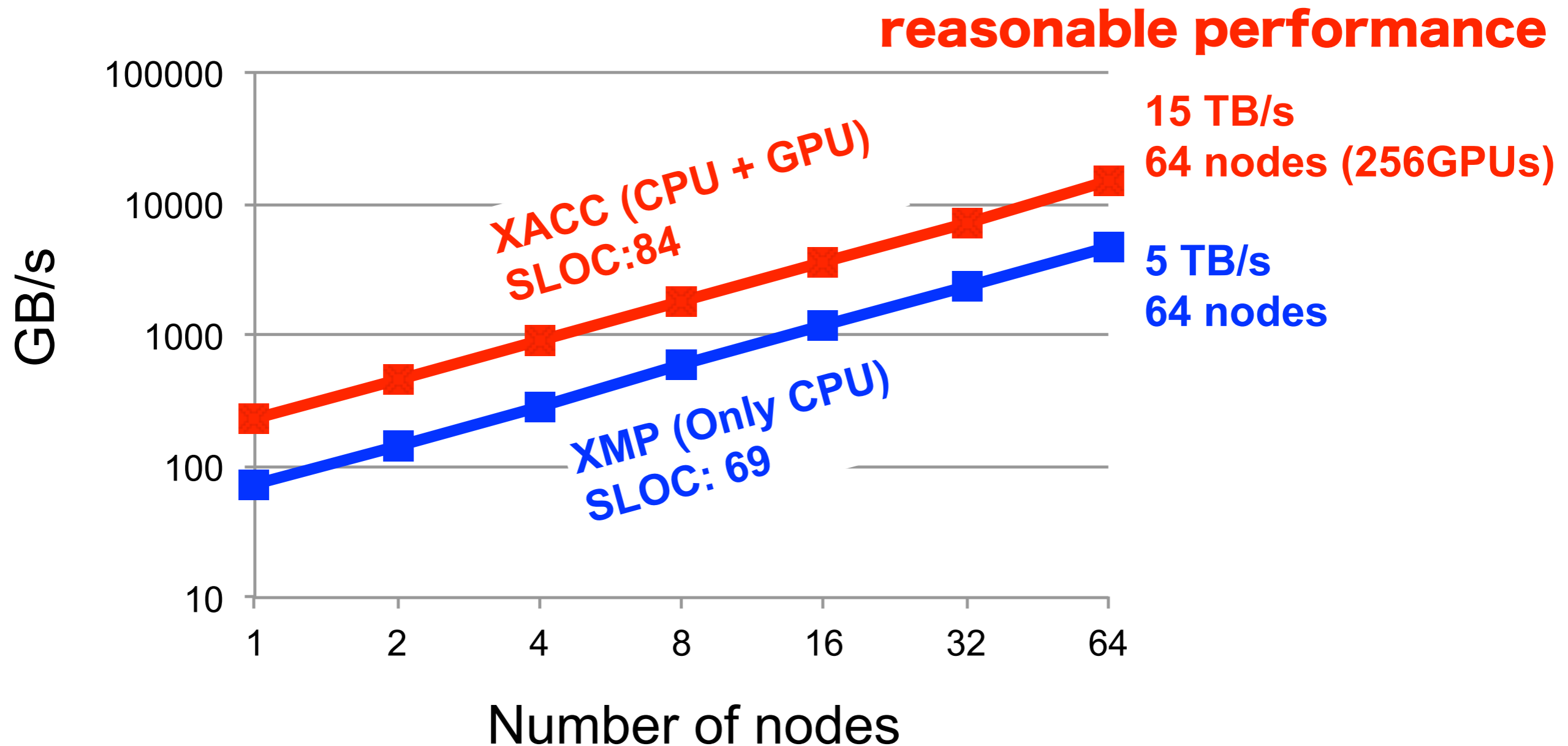
#pragma acc parallel loop async
    for (j=0; j<GPU_SIZE; j++)
      a[j] = b[j] + scalar*c[j];
                                     on GPU

#pragma omp parallel for
    for (j=GPU_SIZE; j<MAX_SIZE; j++)
      a[j] = b[j] + scalar*c[j];
                                     on CPU

#pragma acc wait
                                     Wait until GPU task completes

#pragma xmp barrier
    times[k] += xmp_wtime();
  }
} // acc data
```

# Performance of STREAM



# HIMENO Benchmark

- Stencil application of incompressible fluid analysis code
- Solving the Poisson's equation
- Sequential and MPI Version HIMENO Benchmark is available at

<http://acc.riken.jp/2444.htm>

```
float p[MIMAX][MJMAX][MKMAX];  
// Define distributed array and halo  
  
#pragma acc data copy(p) ..  
{  
..  
#pragma xmp reflect (p) acc  
..  
#pragma xmp loop (k,j,i) on t(k,j,i)  
#pragma acc parallel loop ..  
for(i=1; i<MIMAX; ++i)  
  for(j=1; j<MJMAX; ++j){  
#pragma acc loop vector ..  
  for(k=1; k<MKMAX; ++k){  
    S0 = p[i+1][j][k] * ..;
```

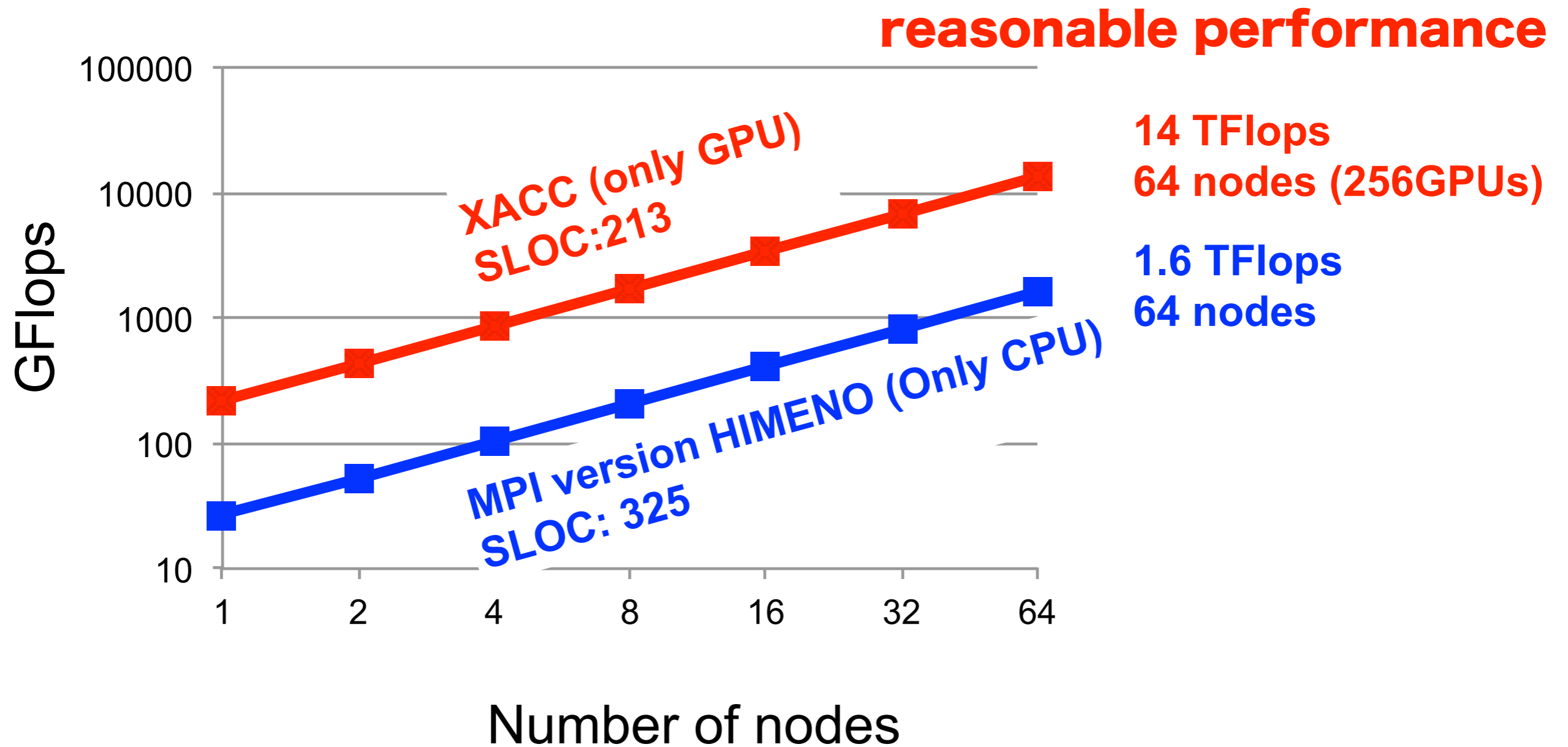
← Transfer distributed array to accelerator

← Exchange halo region

← Parallelize loop statement

Only add **XMP** and **OpenACC** directives into the sequential Himeno benchmark.

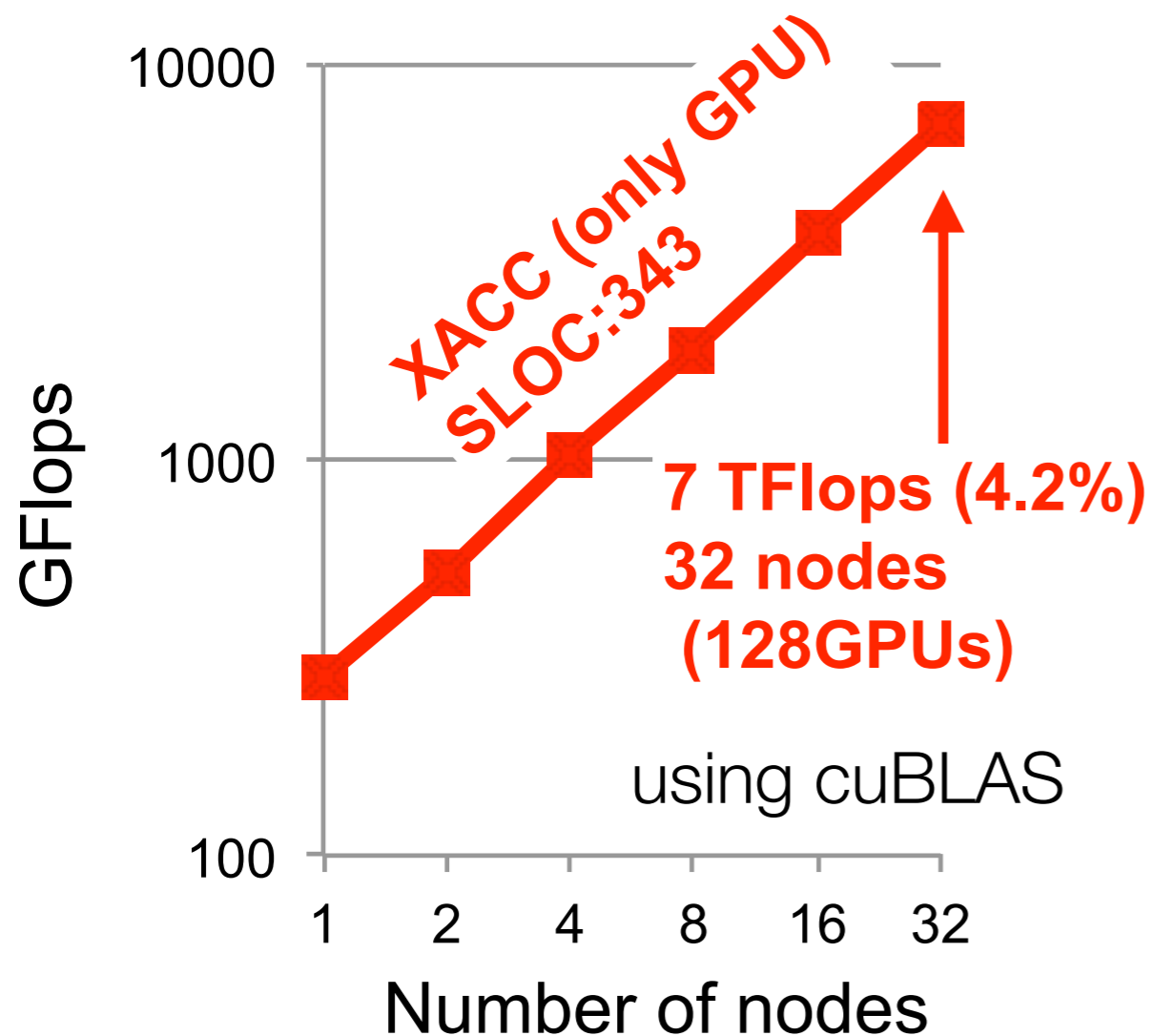
# Performance of HIMENO



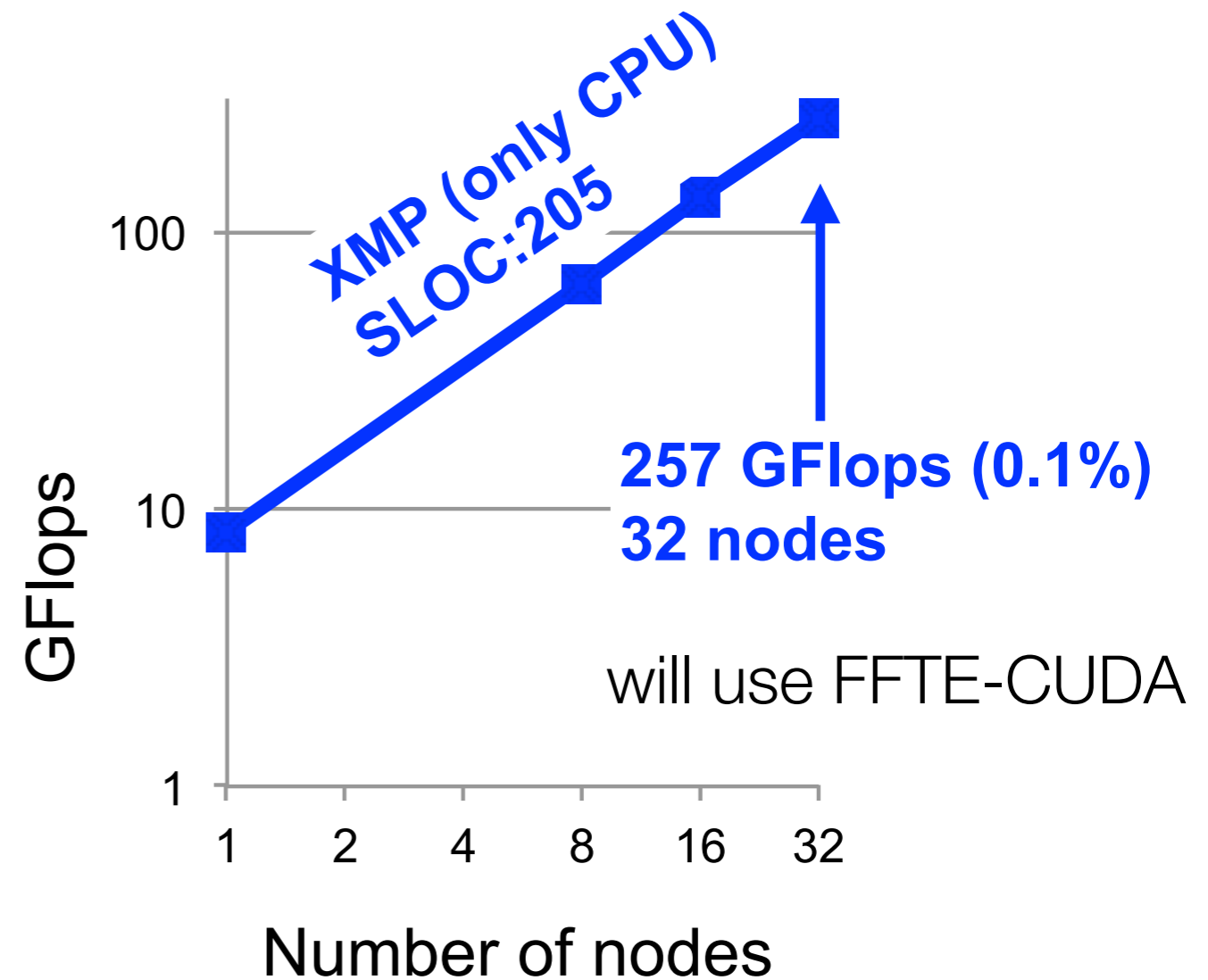
# HPL and FFT

**Sorry !! work-in-progress for implementing and tuning.**

## HPL



## FFT



Time of transfer data between CPU and host memory dominates the total computation time

# Conclusion

## XMP on the K computer

**Good productivity and performance !!**

Benchmark		# Nodes	Performance (/peak)	SLOC
HPL	Ver. 1	16,384	971 TFlops (46.3%)	313
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RandomAccess		16,384	254 GUPs	253

## XACC on HA-PACS/TCA

**We will improve HPL and FFT next year.**

Benchmark	#Nodes	#CPUs	#GPUs	Performance (/peak)	SLOC
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# For more information

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- Center for Computational Sciences, University of Tsukuba #3215